

Canonical & Standard Forms

CSIM601251

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Outline

- **Minterms and Maxterms**
- Relationship between Minterms and Maxterms: The Canonical forms
- Conversion into Canonical Sum-of-Minterm (SOM) or Product-of-Maxterm (POM) Representations
- Standard Form Sum-of-Products (SOP) and Product-of-Sum (POS)

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Minterms and Maxterms (1/2)

- A **minterm** of n variables is a product term that contains n literals from **all** the variables.

Example: On 2 variables x and y , the minterms are:

$$x' \cdot y', x' \cdot y, x \cdot y' \text{ and } x \cdot y$$

- A **maxterm** of n variables is a sum term that contains n literals from **all** the variables.

Example: On 2 variables x and y , the maxterms are:

$$x' + y', x' + y, x + y' \text{ and } x + y$$

- In general, with n variables we have 2^n minterms and 2^n maxterms.

Minterms and Maxterms (2/2)

- The minterms and maxterms on 2 variables are denoted by m_0 to m_3 and M_0 to M_3 respectively.

x	y	Minterms		Maxterms	
		Term	Notation	Term	Notation
0	0	$x' \cdot y'$	m_0	$x+y$	M_0
0	1	$x' \cdot y$	m_1	$x+y'$	M_1
1	0	$x \cdot y'$	m_2	$x'+y$	M_2
1	1	$x \cdot y$	m_3	$x'+y'$	M_3

- Each minterm is the complement of the corresponding maxterm
 - Example: $m_2 = x \cdot y'$
 $m_2' = (x \cdot y')' = x' + (y')' = x' + y = M_2$

Standard Order

- Minterms and maxterms are designated with a subscript
- The subscript is a number, corresponding to a binary pattern
- The bits in the pattern represent the complemented or normal state of each variable listed in a standard order.
- All variables will be present in a minterm or maxterm and will be listed in the same order (usually alphabetically)
- Example: For variables a, b, c :
 - Maxterms: $(a + b + \bar{c}), (a + b + c)$
 - Terms: $(b + a + c), a\bar{c}b$, and $(c + b + a)$ are NOT in standard order.
 - Minterms: $a\bar{b}c, abc, a\bar{b}\bar{c}$
 - Terms: $(a + c), \bar{b}c$, and $(\bar{a} + b)$ do not contain all variables

Purpose of the Index

- The index for the minterm or maxterm, expressed as a binary number, is used to determine whether the variable is shown in the true form or complemented form.
- For Minterms:
 - “1” means the variable is “Not Complemented” and
 - “0” means the variable is “Complemented”.
- For Maxterms:
 - “0” means the variable is “Not Complemented” and
 - “1” means the variable is “Complemented”.

Index Example in Three Variables

- Example: (for three variables)
- Assume the variables are called X, Y, and Z.
- The standard order is X, then Y, then Z.
- The Index 0 (base 10) = 000 (base 2) for three variables).
All three variables are complemented for minterm 0 ($\bar{X}, \bar{Y}, \bar{Z}$) and no variables are complemented for Maxterm 0 (X, Y, Z).
 - Minterm 0, called m_0 is $\bar{X} \cdot \bar{Y} \cdot \bar{Z}$
 - Maxterm 0, called M_0 is $(X + Y + Z)$
 - Minterm 6 ?
 - Maxterm 6 ?

Index Examples - Four Variables

Index	Binary	Minterm	Maxterm
i	Pattern	m_i	M_i
0	0000	$\bar{a}\bar{b}\bar{c}\bar{d}$	$a + b + c + d$
1	0001	$\bar{a}\bar{b}\bar{c}d$?
3	0011	?	$a + b + \bar{c} + \bar{d}$
5	0101	$\bar{a}b\bar{c}d$	$a + \bar{b} + c + \bar{d}$
7	0111	?	$a + \bar{b} + \bar{c} + \bar{d}$
10	1010	$a\bar{b}c\bar{d}$	$\bar{a} + b + \bar{c} + d$
13	1101	$ab\bar{c}d$?
15	1111	$abcd$	$\bar{a} + \bar{b} + \bar{c} + \bar{d}$

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Minterm and Maxterm Relationship

- Review: DeMorgan's Theorem

$$\overline{x \cdot y} = \bar{x} + \bar{y} \text{ and } \overline{x + y} = \bar{x} \cdot \bar{y}$$

- Two-variable example:

$$M_2 = \bar{x} + y \text{ and } m_2 = x \cdot \bar{y}$$

Thus M_2 is the complement of m_2 and vice-versa.

- Since DeMorgan's Theorem holds for n variables, the above holds for terms of n variables
- giving:

$$M_i = \overline{m_i} \text{ and } m_i = \overline{M_i}$$

Thus M_i is the complement of m_i .

Function Tables for Both

- Minterms of 2 variables

x y	m₀	m₁	m₂	m₃
0 0	1	0	0	0
0 1	0	1	0	0
1 0	0	0	1	0
1 1	0	0	0	1

- Maxterms of 2 variables

x y	M₀	M₁	M₂	M₃
0 0	0	1	1	1
0 1	1	0	1	1
1 0	1	1	0	1
1 1	1	1	1	0

- Each column in the maxterm function table is the complement of the column in the minterm function table since M_i is the complement of m_i .

Observations

- In the function tables:
 - Each minterm has one and only one 1 present in the 2^n terms (a minimum of 1s). All other entries are 0.
 - Each maxterm has one and only one 0 present in the 2^n terms All other entries are 1 (a maximum of 1s).
 - We can implement any function by "**ORing**" the minterms corresponding to "1" entries in the function table. These are called the minterms of the function.
 - We can implement any function by "**ANDing**" the maxterms corresponding to "0" entries in the function table. These are called the maxterms of the function.
 - This gives us two canonical forms:
 - Sum of Minterms (SOM)
 - Product of Maxterms (POM)
- for stating any Boolean function.

Minterm Function Example

- Example: Find $F_1 = m_1 + m_4 + m_7$
- $F_1 = \bar{x} \bar{y} z + x \bar{y} \bar{z} + x y z$

x y z	index	m1 + m4 + m7 = F1					
0 0 0	0	0	+	0	+	0	= 0
0 0 1	1	1	+	0	+	0	= 1
0 1 0	2	0	+	0	+	0	= 0
0 1 1	3	0	+	0	+	0	= 0
1 0 0	4	0	+	1	+	0	= 1
1 0 1	5	0	+	0	+	0	= 0
1 1 0	6	0	+	0	+	0	= 0
1 1 1	7	0	+	0	+	1	= 1

Minterm Function Example

- $F(A, B, C, D, E) = m_2 + m_9 + m_{17} + m_{23}$
- $F(A, B, C, D, E) =$

Maxterm Function Example

- Example: Implement F1 in maxterms:

$$F_1 = M_0 \cdot M_2 \cdot M_3 \cdot M_5 \cdot M_6$$

$$F_1 = (x + y + z) \cdot (x + \bar{y} + z) \cdot (x + \bar{y} + \bar{z}) \cdot (\bar{x} + y + \bar{z}) \cdot (\bar{x} + \bar{y} + z)$$

x y z	i	$M_0 \cdot M_2 \cdot M_3 \cdot M_5 \cdot M_6 = F_1$
0 0 0	0	$0 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 0$
0 0 1	1	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 1$
0 1 0	2	$1 \cdot 0 \cdot 1 \cdot 1 \cdot 1 = 0$
0 1 1	3	$1 \cdot 1 \cdot 0 \cdot 1 \cdot 1 = 0$
1 0 0	4	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 1$
1 0 1	5	$1 \cdot 1 \cdot 1 \cdot 0 \cdot 1 = 0$
1 1 0	6	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 0 = 0$
1 1 1	7	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 1$

Maxterm Function Example

- $F(A, B, C, D) = M_3 \times M_8 \times M_{11} \times M_{14}$
- $F(A, B, C, D) =$

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Canonical Sum of Minterms

- Any Boolean function can be expressed as a Sum of Minterms.
 - For the function table, the minterms used are the terms corresponding to the 1's
 - For expressions, expand all terms first to explicitly list all minterms. Do this by “ANDing” any term missing a variable v with a term $(v + \bar{v})$.
- Example: Implement $f = x + \bar{x} \bar{y}$ as a sum of minterms.

First expand terms: $f = x(y + \bar{y}) + \bar{x} \bar{y}$

Then distribute terms: $f = xy + x\bar{y} + \bar{x} \bar{y}$

Express as sum of minterms: $f = m_3 + m_2 + m_0$

Another SOM Example

- Example: $F = A + \bar{B} C$
- There are three variables, A, B, and C which we take to be the standard order.
- Expanding the terms with missing variables:

- Collect terms (removing all but one of duplicate terms):
- Express as SOM:

Shorthand SOM Form

- From the previous example, we started with:

$$F = A + \bar{B} C$$

- We ended up with:

$$F = m_1 + m_4 + m_5 + m_6 + m_7$$

- This can be denoted in the formal shorthand:

$$F(A, B, C) = \Sigma_m(1, 4, 5, 6, 7)$$

- Note that we explicitly show the standard variables in order and drop the “m” designators.

Canonical Product of Maxterms

- Any Boolean Function can be expressed as a Product of Maxterms (POM).
 - For the function table, the maxterms used are the terms corresponding to the 0's.
 - For an expression, expand all terms first to explicitly list all maxterms. Do this by first applying the second distributive law, “ORing” terms missing variable v with a term equal to v and then applying the distributive law again. $V \times \bar{V}$
- Example: Convert to product of maxterms:

$$f(x, y, z) = x + \bar{x} \bar{y}$$

Apply the distributive law:

$$x + \bar{x} \bar{y} = (x + \bar{x})(x + \bar{y}) = 1 \times (x + \bar{y}) = x + \bar{y}$$

Add missing variable z :

$$x + \bar{y} + z \times \bar{z} = (x + \bar{y} + z)(x + \bar{y} + \bar{z})$$

Express as POM: $f = M_2 \cdot M_3$

Another POM Example

- Convert to Product of Maxterms:

$$f(A,B,C) = A\bar{C} + BC + \bar{A}\bar{B}$$

- Use $x + yz = (x+y) \cdot (x+z)$ with $x = (A\bar{C} + BC)$, $y = \bar{A}$, and $z = \bar{B}$ to get:

$$f = (A\bar{C} + BC + \bar{A})(A\bar{C} + BC + \bar{B})$$

- Then use $x + \bar{x}y = x + y$ to get:

$$f = (\bar{C} + BC + \bar{A})(A\bar{C} + C + \bar{B})$$

and a second time to get:

$$f = (\bar{C} + B + \bar{A})(A + C + \bar{B})$$

- Rearrange to standard order,

$$f = (\bar{A} + B + \bar{C})(A + \bar{B} + C) \quad \text{to give } f = M_5 \cdot M_2$$

Another POM Example

Convert to Product of Maxterms: $F_1(A,B,C) = A C' + B C + A'B'$

A	B	C	F_1
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

$$\begin{aligned} F_1 &= M_2 \bullet M_5 \\ &= (A + B' + C) \bullet (A' + B + C') \end{aligned}$$

Function Complements

- The complement of a function expressed as a sum of minterms is constructed by selecting the minterms missing in the sum-of-minterms canonical forms.
- Alternatively, the complement of a function expressed by a Sum of Minterms form is simply the Product of Maxterms with the same indices.
- Example: Given

$$F(x, y, z) = \sum_m(1, 3, 5, 7)$$

$$\bar{F}(x, y, z) = \sum_m(0, 2, 4, 6)$$

$$\bar{F}(x, y, z) = \prod_M(1, 3, 5, 7)$$

Conversion Between Forms

- To convert between sum-of-minterms and product-of-maxterms form (or vice-versa) we follow these steps:
 - Find the function complement by swapping terms in the list with terms not in the list.
 - Change from products to sums, or vice versa.
- Example: Given F as before: $\mathbf{F(x, y, z) = \Sigma_m(1,3,5,7)}$
- Form the Complement: $\mathbf{\bar{F}(x, y, z) = \Sigma_m(0,2,4,6)}$
- Then use the other form with the same indices - this forms the complement again, giving the other form of the original function: $\mathbf{F(x, y, z) = \Pi_M(0,2,4,6)}$

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Standard Forms (1/2)

- Certain types of Boolean expressions lead to circuits that are desirable from implementation viewpoint.
- Two standard forms:
 - Sum-of-Products
 - Product-of-Sums
- Literals
 - A Boolean variable on its own or in its complemented form
 - Examples: x , x' , y , y'
- Product term
 - A single literal or a logical product (AND) of several literals
 - Examples: x , $x \cdot y \cdot z'$, $A' \cdot B$, $A \cdot B$, $d \cdot g' \cdot v \cdot w$

Standard Forms (2/2)

- **Sum term**
 - A single literal or a logical sum (OR) of several literals
 - Examples: x , $x+y+z'$, $A'+B$, $A+B$, $c+d+h'+j$
- **Sum-of-Products (SOP) expression**
 - A product term or a logical sum (OR) of several product terms
 - Examples: x , $x + y \cdot z'$, $x \cdot y' + x' \cdot y \cdot z$, $A \cdot B + A' \cdot B'$,
 $A + B' \cdot C + A \cdot C' + C \cdot D$
- **Product-of-Sums (POS) expression**
 - A sum term or a logical product (AND) of several sum terms
 - Examples: x , $x \cdot (y+z')$, $(x+y') \cdot (x'+y+z)$,
 $(A+B) \cdot (A'+B')$, $(A+B+C) \cdot D' \cdot (B'+D+E')$
- **Every Boolean expression can be expressed in SOP or POS.**

Do it yourself

- Put the right ticks in the following table.

<i>Expression</i>	<i>SOP?</i>	<i>POS?</i>
$X' \cdot Y + X \cdot Y' + X \cdot Y \cdot Z$		
$(X+Y') \cdot (X'+Y) \cdot (X'+Z')$		
$X' + Y + Z$		
$X \cdot (W' + Y \cdot Z)$		
$X \cdot Y \cdot Z'$		
$W \cdot X' \cdot Y + V \cdot (X \cdot Z + W')$		

Standard Sum-of-Products (SOP)

- A sum of minterms form for n variables can be written down directly from a truth table.
 - Implementation of this form is a two-level network of gates such that:
 - The first level consists of n -input AND gates, and
 - The second level is a single OR gate (with fewer than 2^n inputs).
- This form often can be simplified so that the corresponding circuit is simpler.

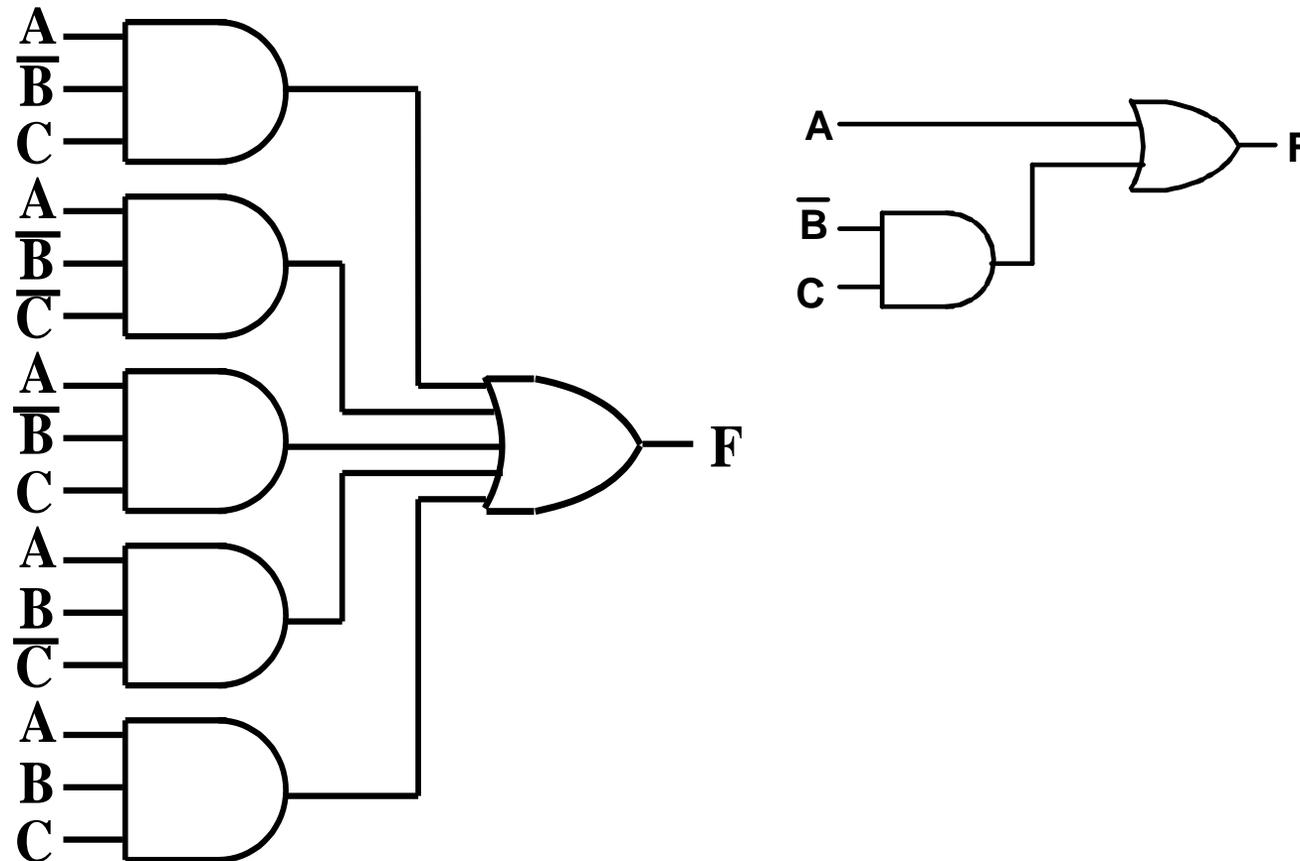
Standard Sum-of-Products (SOP)

- A Simplification Example:
- $F(A, B, C) = \Sigma m(1,4,5,6,7)$
- Writing the minterm expression:
$$F = \overline{A} \overline{B} C + A \overline{B} \overline{C} + A \overline{B} C + ABC\overline{C} + ABC$$
- Simplifying:
$$F =$$

- Simplified F contains 3 literals compared to 15 in minterm F

AND/OR Two-level Implementation of SOP Expression

- The two implementations for F are shown below – it is quite apparent which is simpler!



SOP and POS Observations

- The previous examples show that:
 - Canonical Forms (Sum-of-minterms, Product-of-Maxterms), or other standard forms (SOP, POS) differ in complexity
 - Boolean algebra can be used to manipulate equations into simpler forms.
 - Simpler equations lead to simpler two-level implementations
- Questions:
 - How can we attain a “simplest” expression?
 - Is there only one minimum cost circuit?
 - The next part will deal with these issues.