

Pemrograman Fungsional

Komputasi, Ekspresi, Datatype



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FAKULTAS
ILMU
KOMPUTER

Review Diskusi Pekan 01 Scele

- Bahasa lain dalam rumpun deklaratif

Prolog

SQL



Review Diskusi Pekan 01 Scele

- Bahasa lain dalam paradigma fungsional
 - Erlang
 - Elixir
 - Miranda
 - Lisp
 - Clojure
 - F#
 - Javascript



Review Diskusi Pekan 01 Scele

- GraphQL – SQL, declarative?
 - GraphQL termasuk declarative seperti SQL
 - Belum akan sepenuhnya menggantikan microservices



Review Diskusi Pekan 01 Scele

- Lambda di Java
 - Anonymous Inner Class
 - Stream API



Review Diskusi Pekan 01 Scele

- Goodbye OO?
 - No silver bullet
 - Masing-masing punya kelebihan
 - Strukturisasi dan abstraksi dari OO masih membantu
 - Selalu kritis!



Agenda dan Learning Objective

- Agenda
 - Model komputasi
 - Ekspresi, values, type
 - Datatype
- Learning Objective
 - Memahami model komputasi Fungsional. Diberikan sebuah program, dapat menjelaskan cara proses eksekusinya
 - Memahami konsep ekspresi sebagai bagian dari proses pembuatan algoritma. Diberikan permasalahan, dapat menyusun ekspresi sebagai penyusun algoritma dan penyelesaian masalah
 - Memahami representasi datatype. Diberikan permodelan masalah, dapat menyusun abstract data type yang berkesesuaian.



Computation by Calculation

- Computation by calculation is a simple concept that everyone should be familiar with, since it's not unlike ordinary arithmetic calculation. For example:

$$\begin{aligned} &3 * (9 + 5) \\ &\rightarrow 3 * 14 \\ &\rightarrow 42 \end{aligned}$$

- However, since we want computers to perform these tasks, we are also interested in issues such as efficiency:

$$\begin{aligned} &3 * (9 + 5) \\ &\rightarrow 3*9 + 3*5 \\ &\rightarrow 27 + 3*5 \\ &\rightarrow 27 + 15 \\ &\rightarrow 42 \end{aligned}$$

- Same answer, but the former was *more efficient* than the latter, since it took a fewer number of steps.



Abstraction

- We are also interested in *abstraction*: the process of recognizing a repeating pattern, and capturing it succinctly in one place instead of many.

- For example:

$$3 * (9 + 5) \quad 4 * (6 + 2) \quad 7 * (8 + 1) \quad \dots$$

This repeating pattern can be captured as a *function*:

$$\text{easy } x \ y \ z = x * (y + z)$$

Then each instance can be replaced by:

$$\text{easy } 3 \ 9 \ 5 \quad \text{easy } 4 \ 6 \ 2 \quad \text{easy } 7 \ 8 \ 1 \quad \dots$$

- We can also perform calculations with *symbols*. For example, we can *prove* that $\text{easy } a \ b \ c = \text{easy } a \ c \ b$:

$$\text{easy } a \ b \ c$$

$$\rightarrow a * (b + c) \quad \{ \text{unfold} \}$$

$$\rightarrow a * (c + b) \quad \{ \text{commutativity of } + \}$$

$$\rightarrow \text{easy } a \ c \ b \quad \{ \text{fold} \}$$



Expressions, Values, and Types

- The objects on which we calculate are called *expressions*.
- When no more unfolding (of either a primitive or user-defined function) is possible, the resulting expression is called a *value*.
- Every expression (and therefore every value) has a *type*.
(A type is a collection of expressions with common attributes.)
- We write $\mathbf{exp} :: \mathbf{T}$ to say that expression \mathbf{exp} has type \mathbf{T} .
- Examples:
 - Atomic expressions:
 $42 :: \mathbf{Integer}, \quad 'a' :: \mathbf{Char}, \quad \mathbf{True} :: \mathbf{Bool}$
 - Structured expressions:
 $[1,2,3] :: [\mathbf{Integer}]$ - a *list* of integers
 $('b', 4) :: (\mathbf{Char}, \mathbf{Integer})$ - a *pair* consisting of a character and an integer
 - Functions:
 $(+) :: \mathbf{Integer} \rightarrow \mathbf{Integer} \rightarrow \mathbf{Integer}$
 $\mathbf{easy} :: \mathbf{Integer} \rightarrow \mathbf{Integer} \rightarrow \mathbf{Integer} \rightarrow \mathbf{Integer}$



Abstraction

- Our derivation of the function **easy** is a good example of the use of the *abstraction principle*: separating a repeating pattern from the particular instances in which it appears. In particular, the example demonstrates *functional abstraction*.
- *Naming* is an even simpler kind of abstraction:

```
let pi = 3.14159
in 2*pi*r1 + 2*pi*r2
```
- *Data abstraction* is the use of data structures to store common values on which common operations may be applied in an abstract manner.
- The “circle areas” example from the text demonstrates all three kinds of abstraction.



“Circle Areas” Example

- Original program:

```
totalArea = pi*r1^2 + pi*r2^2 + pi*r3^2
```

- A more abstract program:

```
totalArea =
```

```
  listSum [circArea r1, circArea r2, circArea r3]
```

```
listSum [] = 0
```

```
listSum (a:as) = a + listSum as
```

```
circArea r = pi*r^2
```

- The new program is longer than the old – in what ways is it better?
 - The code for the area of a circle has been isolated (using functional abstraction), thus minimizing errors and facilitating change and reuse.
 - The number of circles has been generalized (via data abstraction) from three to an arbitrary number, thus anticipating change and reuse.



Evaluation: Proof by Calculation

- Proof that the new `totalArea` is equivalent to the old:

```
listSum [circArea r1, circArea r2, circArea r3]
```

```
→ { unfold listSum }
```

```
circArea r1 + listSum [circArea r2, circArea r3]
```

```
→ { unfold listSum }
```

```
circArea r1 + circArea r2 + listSum [circArea r3]
```

```
→ { unfold listSum }
```

```
circArea r1 + circArea r2 + circArea r3 + listSum []
```

```
→ { unfold listSum }
```

```
circArea r1 + circArea r2 + circArea r3 + 0
```

```
→ { unfold circArea (three places) }
```

```
pi*r1^2 + pi*r2^2 + pi*r3^2 + 0
```

```
→ { simple arithmetic }
```

```
pi*r1^2 + pi*r2^2 + pi*r3^2
```



Defining New Datatypes

- The ability to define new data types in a programming language is important.
- Kinds of data types:
 - *enumerated types*
 - *records (or products)*
 - *variant records (or sums)*
 - *recursive types*
- Haskell's `data` declaration provides these kinds of data types in a uniform way that abstracts away from their implementation details, by providing an abstract interface to the newly defined type.
- Before looking at the example from Chapter 2, let's look at some simpler examples.



The Data Declaration

- Example of an *enumeration* data type:

```
data Day = Sun | Mon | Tue | Wed | Thu | Fri | Sat
  deriving Show
```

- The names **Sun through Sat** are constructor constants (since they have no arguments) and are the *only* elements of the type.
- For example, we can define:

```
valday :: Integer -> Day
valday 1 = Sun
valday 2 = Mon
valday 3 = Tue
valday 4 = Wed
valday 5 = Thu
valday 6 = Fri
valday 7 = Sat
```

```
Hugs> valday 4
Wed
```



Constructors & Patterns

- Constructors can be matched by *patterns*.
- For example :

```
dayval :: Day -> Integer
dayval Sun = 1
dayval Mon = 2
dayval Tue = 3
dayval Wed = 4
dayval Thu = 5
dayval Fri = 6
dayval Sat = 7
```

```
Hugs> dayval Wed
4
```



Other Enumeration Data Type Examples

```
data Bool = True | False -- predefined in Haskell
    deriving Show
data Direction = North | East | South | West
    deriving Show
data Move = Paper | Rock | Scissors
    deriving Show
```

```
beats :: Move -> Move
beats Paper      = Scissors
beats Rock       = Paper
beats Scissors   = Rock
```

```
Hugs> beats Paper
Scissors
```



Variant Records

- More complicated data types:

```
data Tagger = Tagn Integer | Tagb Bool
```

- These constructors are not constants – they are **functions**:

```
Tagn :: Integer -> Tagger
```

```
Tagb :: Bool -> Tagger
```

- As for all constructors, something like “**Tagn 12**”
 - Cannot be simplified
(and thus, as discussed in Chapter 1, it is a *value*).
 - Can be used in patterns.



Example functions on **Tagger**

```
number (Tagn n) = n
```

```
boolean (Tagb b) = b
```

```
isNum (Tagn _) = True
```

```
isNum (Tagb _) = False
```

```
isBool x = not (isNum x)
```

```
Hugs> :t number
```

```
number :: Tagger -> Integer
```

```
Hugs> number (Tagn 3)
```

```
3
```

```
Hugs> isNum (Tagb False)
```

```
False
```



Another Variant Record Data Type

```
data Temp = Celsius Float
          | Fahrenheit Float
          | Kelvin Float
```

- We can use patterns to define functions over this type:

```
toKelvin (Celsius c)      = Kelvin (c + 272.0)
toKelvin (Fahrenheit f) =
    Kelvin ( 5/9*(f-32.0) + 272.0 )
toKelvin (Kelvin k)      = Kelvin k
```



Finally: the **Shape** Data Type from the Text

- The **Shape** data type from Chapter 2 is another example of a variant data type:

```
data Shape = Rectangle Float Float
           | Ellipse Float Float
           | RtTriangle Float Float
           | Polygon [(Float,Float)]
deriving Show
```

- The last line – “**deriving Show**” – tells the system to build a **show** function for the type **Shape** (more on this later).
- We can also define functions yielding refined shapes:

```
circle, square :: Float -> Shape
circle radius = Ellipse radius radius
square side   = Rectangle side side
```



Functions over `shape`

- Functions on shapes can be defined using pattern matching.

```
area :: Shape -> Float
area (Rectangle s1 s2) = s1*s2
area (Ellipse r1 r2)   = pi*r1*r2
area (RtTriangle s1 s2) = (s1*s2)/2
area (Polygon (v1:pts)) = polyArea pts
  where polyArea :: [(Float,Float)] -> Float
        polyArea (v2:v3:vs) = triArea v1 v2 v3 +
                               polyArea (v3:vs)
        polyArea _          = 0
```

Note use of auxiliary function.

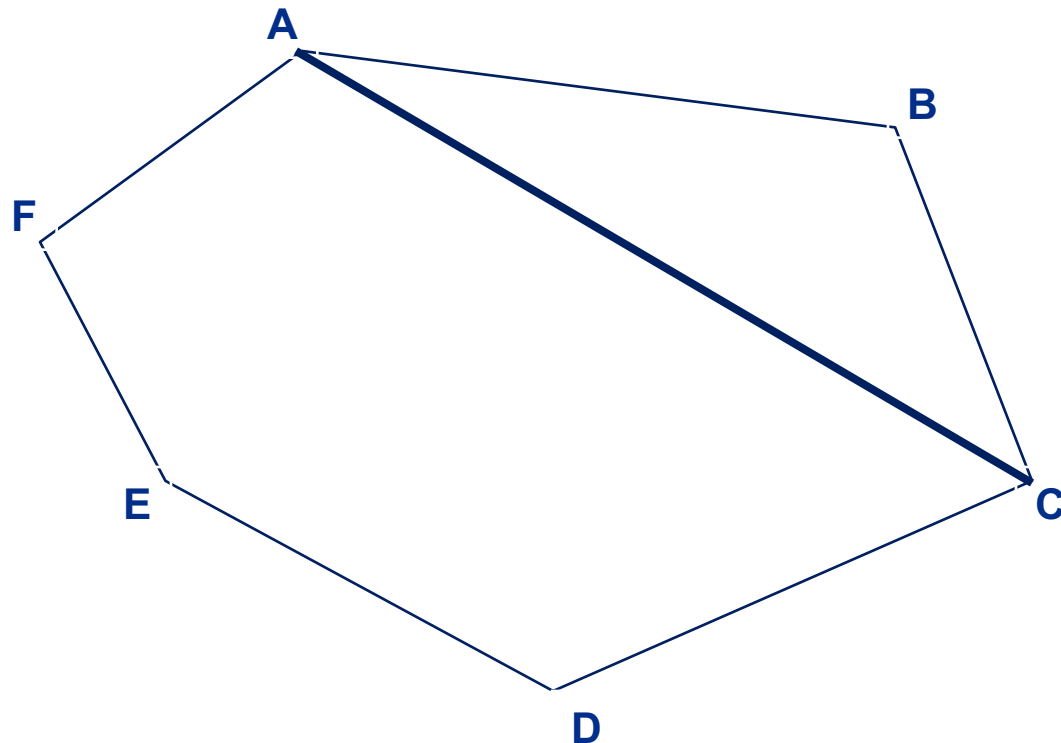
Note use of nested patterns.

Note use of wild card pattern (which matches anything).



Algorithm for Computing Area of Polygon

$\text{totalArea} = \text{area}(\text{triangle}[A,B,C]) + \text{area}(\text{polygon}[A,C,D,E,F])$



TriArea

```
triArea v1 v2 v3 =  
  let a = distBetween v1 v2  
      b = distBetween v2 v3  
      c = distBetween v3 v1  
      s = 0.5*(a+b+c)  
  in sqrt (s*(s-a)*(s-b)*(s-c))
```

```
distBetween (x1,y1) (x2,y2)  
  = sqrt ((x1-x2)^2 + (y1-y2)^2)
```





Selamat Belajar dan Berlatih!

Silahkan baca bab 1-3 buku Haskell School of
Expression

atau bab 1-4 buku Haskell The Craft of Functional
Programming



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