

# Canonical & Standard Forms

CSIM601251

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# Outline

- **Minterms and Maxterms**
- Relationship between Minterms and Maxterms: The Canonical forms
- Conversion into Canonical Sum-of-Minterm (SOM) or Product-of-Maxterm (POM) Representations
- Standard Form Sum-of-Products (SOP) and Product-of-Sum (POS)

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# Minterms and Maxterms (1/2)

- A **minterm** of  $n$  variables is a product term that contains  $n$  literals from **all** the variables.

Example: On 2 variables  $x$  and  $y$ , the minterms are:

$$x' \cdot y', x' \cdot y, x \cdot y' \text{ and } x \cdot y$$

- A **maxterm** of  $n$  variables is a sum term that contains  $n$  literals from **all** the variables.

Example: On 2 variables  $x$  and  $y$ , the maxterms are:

$$x' + y', x' + y, x + y' \text{ and } x + y$$

- In general, with  $n$  variables we have  $2^n$  minterms and  $2^n$  maxterms.

# Minterms and Maxterms (2/2)

- The minterms and maxterms on 2 variables are denoted by **m0 to m3** and **M0 to M3** respectively.

x	y	Minterms		Maxterms	
		Term	Notation	Term	Notation
0	0	$x' \cdot y'$	m0	$x+y$	M0
0	1	$x' \cdot y$	m1	$x+y'$	M1
1	0	$x \cdot y'$	m2	$x'+y$	M2
1	1	$x \cdot y$	m3	$x'+y'$	M3

- Each minterm is the complement of the corresponding maxterm
  - Example:  $m2 = x \cdot y'$   
 $m2' = (x \cdot y')' = x' + (y')' = x' + y = M2$

# Standard Order

- Minterms and maxterms are designated with a subscript
- The subscript is a number, corresponding to a binary pattern
- The bits in the pattern represent the complemented or normal state of each variable listed in a standard order.
- All variables will be present in a minterm or maxterm and will be listed in the same order (usually alphabetically)
- Example: For variables a, b, c:
  - Maxterms:  $(a + b + \bar{c})$ ,  $(a + b + c)$
  - Terms:  $(b + a + c)$ ,  $a\bar{c}b$ , and  $(c + b + a)$  are NOT in standard order.
  - Minterms:  $a\bar{b}c$ ,  $abc$ ,  $a\bar{b}\bar{c}$
  - Terms:  $(a + c)$ ,  $\bar{b}c$ , and  $(\bar{a} + b)$  do not contain all variables

# Purpose of the Index

- The index for the minterm or maxterm, expressed as a binary number, is used to determine whether the variable is shown in the true form or complemented form.
- For Minterms:
  - “1” means the variable is “Not Complemented” and
  - “0” means the variable is “Complemented”.
- For Maxterms:
  - “0” means the variable is “Not Complemented” and
  - “1” means the variable is “Complemented”.

# Index Example in Three Variables

- Example: (for three variables)
- Assume the variables are called X, Y, and Z.
- The standard order is X, then Y, then Z.
- The Index 0 (base 10) = 000 (base 2) for three variables).  
All three variables are complemented for minterm 0  
( $\bar{X}, \bar{Y}, \bar{Z}$ ) and no variables are complemented for Maxterm 0  
0 (X, Y, Z).
  - Minterm 0, called  $m_0$  is  $\bar{X} \cdot \bar{Y} \cdot \bar{Z}$
  - Maxterm 0, called  $M_0$  is  $(X + Y + Z)$
  - Minterm 6 ?
  - Maxterm 6 ?

# Index Examples – Four Variables

Index	Binary	Minterm	Maxterm
i	Pattern	$m_i$	$M_i$
0	0000	$\bar{a}\bar{b}\bar{c}\bar{d}$	$a + b + c + d$
1	0001	$\bar{a}\bar{b}\bar{c}d$	?
3	0011	?	$a + b + \bar{c} + \bar{d}$
5	0101	$\bar{a}b\bar{c}d$	$a + \bar{b} + c + \bar{d}$
7	0111	?	$a + \bar{b} + \bar{c} + \bar{d}$
10	1010	$a\bar{b}c\bar{d}$	$\bar{a} + b + \bar{c} + d$
13	1101	$ab\bar{c}d$	?
15	1111	$abcd$	$\bar{a} + \bar{b} + \bar{c} + \bar{d}$



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- Conversion into Canonical Sum-of-Minterm (SOM) or Product-of-Maxterm (POM) Representations
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# Minterm and Maxterm Relationship

- Review: DeMorgan's Theorem

$$\overline{x \cdot y} = \bar{x} + \bar{y} \text{ and } \overline{x + y} = \bar{x} \cdot \bar{y}$$

- Two-variable example:

$$M_2 = \bar{x} + y \text{ and } m_2 = x \cdot \bar{y}$$

Thus  $M_2$  is the complement of  $m_2$  and vice-versa.

- Since DeMorgan's Theorem holds for  $n$  variables, the above holds for terms of  $n$  variables
- giving:

$$M_i = \overline{m_i} \text{ and } m_i = \overline{M_i}$$

Thus  $M_i$  is the complement of  $m_i$ .

# Function Tables for Both

- Minterms of 2 variables

x y	$m_0$	$m_1$	$m_2$	$m_3$
0 0	1	0	0	0
0 1	0	1	0	0
1 0	0	0	1	0
1 1	0	0	0	1

- Maxterms of 2 variables

x y	$M_0$	$M_1$	$M_2$	$M_3$
0 0	0	1	1	1
0 1	1	0	1	1
1 0	1	1	0	1
1 1	1	1	1	0

- Each column in the maxterm function table is the complement of the column in the minterm function table since  $M_i$  is the complement of  $m_i$ .

# Observations

- In the function tables:
  - Each minterm has one and only one 1 present in the  $2^n$  terms (a minimum of 1s). All other entries are 0.
  - Each maxterm has one and only one 0 present in the  $2^n$  terms All other entries are 1 (a maximum of 1s).
- We can implement any function by "**OR**ing" the minterms corresponding to "1" entries in the function table. These are called the minterms of the function.
- We can implement any function by "**AND**ing" the maxterms corresponding to "0" entries in the function table. These are called the maxterms of the function.
- This gives us two canonical forms:
  - Sum of Minterms (SOM)
  - Product of Maxterms (POM)

for stating any Boolean function.

# Minterm Function Example

- Example: Find  $F_1 = m_1 + m_4 + m_7$
- $F_1 = \bar{x} \bar{y} z + x \bar{y} \bar{z} + x y z$

x y z	index	m1 + m4 + m7 = F1				
0 0 0	0	0	+	0	+	0 = 0
0 0 1	1	1	+	0	+	0 = 1
0 1 0	2	0	+	0	+	0 = 0
0 1 1	3	0	+	0	+	0 = 0
1 0 0	4	0	+	1	+	0 = 1
1 0 1	5	0	+	0	+	0 = 0
1 1 0	6	0	+	0	+	0 = 0
1 1 1	7	0	+	0	+	1 = 1

# Minterm Function Example

- $F(A, B, C, D, E) = m_2 + m_9 + m_{17} + m_{23}$
- $F(A, B, C, D, E) =$

# Maxterm Function Example

- Example: Implement F1 in maxterms:

$$F_1 = M_0 \cdot M_2 \cdot M_3 \cdot M_5 \cdot M_6$$

$$F_1 = (x + y + z) \cdot (x + \bar{y} + z) \cdot (x + \bar{y} + \bar{z}) \cdot (\bar{x} + y + \bar{z}) \cdot (\bar{x} + \bar{y} + z)$$

x y z	i	$M_0 \cdot M_2 \cdot M_3 \cdot M_5 \cdot M_6 = F1$
0 0 0	0	$0 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 0$
0 0 1	1	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 1$
0 1 0	2	$1 \cdot 0 \cdot 1 \cdot 1 \cdot 1 = 0$
0 1 1	3	$1 \cdot 1 \cdot 0 \cdot 1 \cdot 1 = 0$
1 0 0	4	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 1$
1 0 1	5	$1 \cdot 1 \cdot 1 \cdot 0 \cdot 1 = 0$
1 1 0	6	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 0 = 0$
1 1 1	7	$1 \cdot 1 \cdot 1 \cdot 1 \cdot 1 = 1$

# Maxterm Function Example

- $F(A, B, C, D) = M_3 \times M_8 \times M_{11} \times M_{14}$
- $F(A, B, C, D) =$



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# Canonical Sum of Minterms

- Any Boolean function can be expressed as a Sum of Minterms.
  - For the function table, the minterms used are the terms corresponding to the 1's
  - For expressions, expand all terms first to explicitly list all minterms. Do this by “ANDing” any term missing a variable  $v$  with a term  $(v + \bar{v})$ .
- Example: Implement  $f = x + \bar{x} \bar{y}$  as a sum of minterms.

First expand terms:  $f = x(y + \bar{y}) + \bar{x} \bar{y}$

Then distribute terms:  $f = xy + x\bar{y} + \bar{x} \bar{y}$

Express as sum of minterms:  $f = m_3 + m_2 + m_0$

# Another SOM Example

- Example:  $F = A + \bar{B} C$
- There are three variables, A, B, and C which we take to be the standard order.
- Expanding the terms with missing variables:
- Collect terms (removing all but one of duplicate terms):
- Express as SOM:

# Shorthand SOM Form

- From the previous example, we started with:

$$F = A + \bar{B} C$$

- We ended up with:

$$F = m_1 + m_4 + m_5 + m_6 + m_7$$

- This can be denoted in the formal shorthand:

$$F(A, B, C) = \Sigma_m(1, 4, 5, 6, 7)$$

- Note that we explicitly show the standard variables in order and drop the “m” designators.

# Canonical Product of Maxterms

- Any Boolean Function can be expressed as a Product of Maxterms (POM).
  - For the function table, the maxterms used are the terms corresponding to the 0's.
  - For an expression, expand all terms first to explicitly list all maxterms. Do this by first applying the second distributive law, "ORing" terms missing variable  $v$  with a term equal to  $v \times \bar{v}$  and then applying the distributive law again.  $\mathbf{V \times \bar{V}}$
- Example: Convert to product of maxterms:

$$f(x, y, z) = x + \bar{x} \bar{y}$$

Apply the distributive law:

$$x + \bar{x} \bar{y} = (x + \bar{x})(x + \bar{y}) = 1 \times (x + \bar{y}) = x + \bar{y}$$

Add missing variable  $z$ :

$$x + \bar{y} + z \times \bar{z} = (x + \bar{y} + z)(x + \bar{y} + \bar{z})$$

Express as POM:  $f = M_2 \cdot M_3$

# Another POM Example

- Convert to Product of Maxterms:

$$f(A,B,C) = A\bar{C} + BC + \bar{A}\bar{B}$$

- Use  $x + yz = (x+y) \cdot (x+z)$  with  $x = (A\bar{C} + BC)$ ,  $y = \bar{A}$ , and  $z = \bar{B}$  to get:

$$f = (A\bar{C} + BC + \bar{A})(A\bar{C} + BC + \bar{B})$$

- Then use  $x + \bar{x}y = x + y$  to get:

$$f = (\bar{C} + BC + \bar{A})(A\bar{C} + C + \bar{B})$$

and a second time to get:

$$f = (\bar{C} + B + \bar{A})(A + C + \bar{B})$$

- Rearrange to standard order,

$$f = (\bar{A} + B + \bar{C})(A + \bar{B} + C) \text{ to give } f = M_5 \cdot M_2$$

# Another POM Example

Convert to Product of Maxterms:  $F_1(A,B,C) = A C' + B C + A' B'$

A	B	C	$F_1$
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

$$\begin{aligned} F_1 &= M_2 \bullet M_5 \\ &= (A + B' + C) \bullet (A' + B + C') \end{aligned}$$

# Function Complements

- The complement of a function expressed as a sum of minterms is constructed by selecting the minterms missing in the sum-of-minterms canonical forms.
- Alternatively, the complement of a function expressed by a Sum of Minterms form is simply the Product of Maxterms with the same indices.
- Example: Given

$$F(x, y, z) = \sum_m(1, 3, 5, 7)$$

$$\bar{F}(x, y, z) = \sum_m(0, 2, 4, 6)$$

$$\bar{F}(x, y, z) = \prod_M(1, 3, 5, 7)$$



# Conversion Between Forms

- To convert between sum-of-minterms and product-of-maxterms form (or vice-versa) we follow these steps:
  - Find the function complement by swapping terms in the list with terms not in the list.
  - Change from products to sums, or vice versa.
- Example: Given  $F$  as before:  $F(x, y, z) = \sum_m(1, 3, 5, 7)$
- Form the Complement:  $\bar{F}(x, y, z) = \sum_m(0, 2, 4, 6)$
- Then use the other form with the same indices – this forms the complement again, giving the other form of the original function:  
 $F(x, y, z) = \prod_M(0, 2, 4, 6)$

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# Standard Forms (1/2)

- Certain types of Boolean expressions lead to circuits that are desirable from implementation viewpoint.
- Two standard forms:
  - Sum-of-Products
  - Product-of-Sums
- Literals
  - A Boolean variable on its own or in its complemented form
  - Examples:  $x$ ,  $x'$ ,  $y$ ,  $y'$
- Product term
  - A single literal or a logical product (AND) of several literals
  - Examples:  $x$ ,  $x \cdot y \cdot z'$ ,  $A' \cdot B$ ,  $A \cdot B$ ,  $d \cdot g' \cdot v \cdot w$

# Standard Forms (2/2)

- **Sum term**
  - A single literal or a logical sum (OR) of several literals
  - Examples:  $x$ ,  $x+y+z'$ ,  $A'+B$ ,  $A+B$ ,  $c+d+h'+j$
- **Sum-of-Products (SOP) expression**
  - A product term or a logical sum (OR) of several product terms
  - Examples:  $x$ ,  $x + y \cdot z'$ ,  $x \cdot y' + x' \cdot y \cdot z$ ,  $A \cdot B + A' \cdot B'$ ,  
 $A + B' \cdot C + A \cdot C' + C \cdot D$
- **Product-of-Sums (POS) expression**
  - A sum term or a logical product (AND) of several sum terms
  - Examples:  $x$ ,  $x \cdot (y+z')$ ,  $(x+y') \cdot (x'+y+z)$ ,  
 $(A+B) \cdot (A'+B')$ ,  $(A+B+C) \cdot D' \cdot (B'+D+E')$
- Every Boolean expression can be expressed in SOP or POS.

# Do it yourself

- Put the right ticks in the following table.

<i>Expression</i>	<i>SOP?</i>	<i>POS?</i>
$X' \cdot Y + X \cdot Y' + X \cdot Y \cdot Z$		
$(X+Y') \cdot (X'+Y) \cdot (X'+Z')$		
$X' + Y + Z$		
$X \cdot (W' + Y \cdot Z)$		
$X \cdot Y \cdot Z'$		
$W \cdot X' \cdot Y + V \cdot (X \cdot Z + W')$		

# Standard Sum-of-Products (SOP)

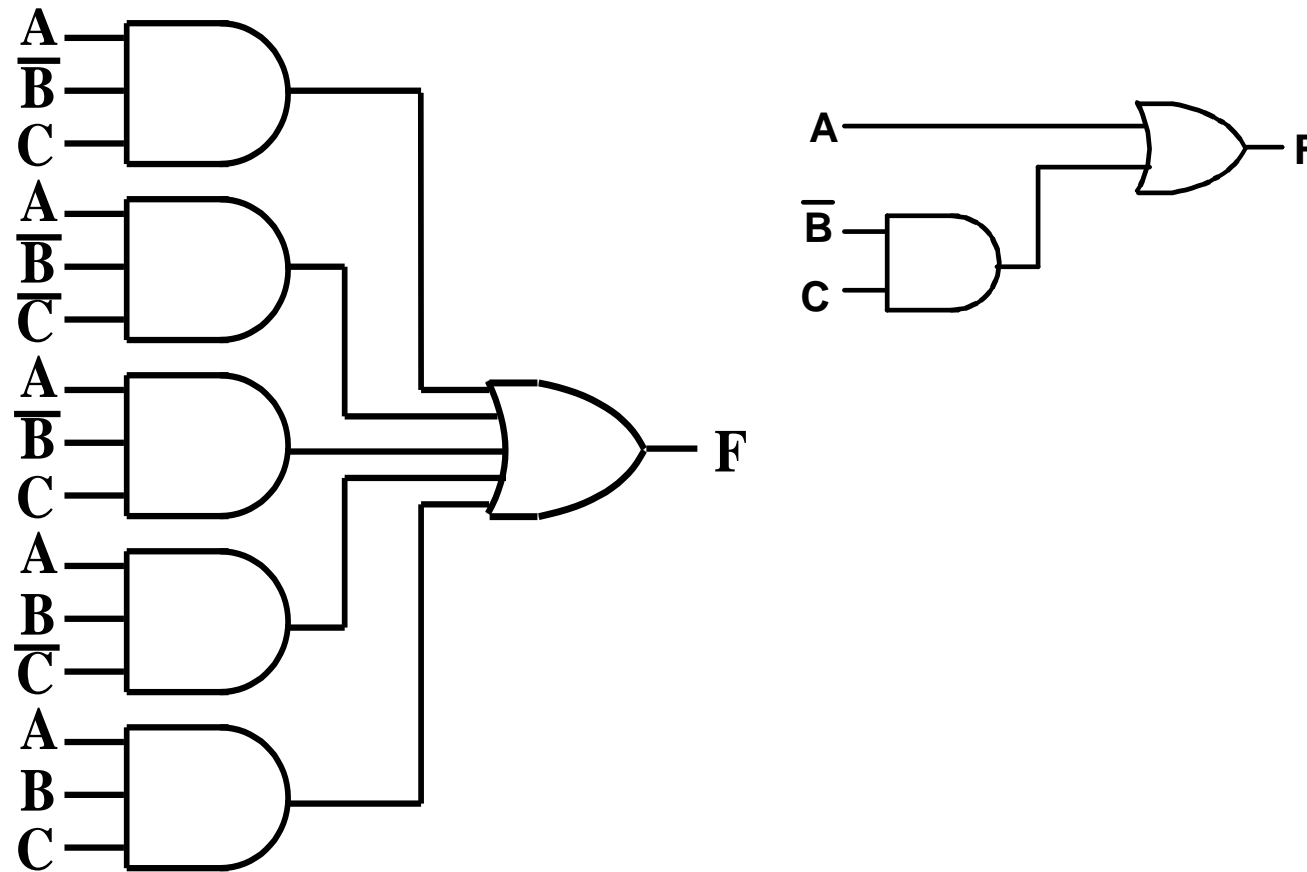
- A sum of minterms form for  $n$  variables can be written down directly from a truth table.
  - Implementation of this form is a two-level network of gates such that:
    - The first level consists of  $n$ -input AND gates, and
    - The second level is a single OR gate (with fewer than  $2^n$  inputs).
- This form often can be simplified so that the corresponding circuit is simpler.

# Standard Sum-of-Products (SOP)

- A Simplification Example:
- $F(A, B, C) = \Sigma m(1, 4, 5, 6, 7)$
- Writing the minterm expression:  
$$F = \overline{A} \overline{B} C + A \overline{B} \overline{C} + A \overline{B} C + A B \overline{C} + A B C$$
- Simplifying:  
$$F =$$
- Simplified F contains 3 literals compared to 15 in minterm F

# AND/OR Two-level Implementation of SOP Expression

- The two implementations for F are shown below – it is quite apparent which is simpler!





# SOP and POS Observations

- The previous examples show that:
  - Canonical Forms (Sum-of-minterms, Product-of-Maxterms), or other standard forms (SOP, POS) differ in complexity
  - Boolean algebra can be used to manipulate equations into simpler forms.
  - Simpler equations lead to simpler two-level implementations
- Questions:
  - How can we attain a “simplest” expression?
  - Is there only one minimum cost circuit?
  - The next part will deal with these issues.